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## Minimum projection uniformity for computer experiment with quantitative factors

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## ABSTRACT

Computer experiments involving quantitative factors at high levels are becoming more and more important in the study of complex experiments arising in the area of science and engineering. Uniform designs are found to be widely applicable in computer experiments in the form of space-filling designs. In this paper, the projection uniformity for quantitative designs is studied under wrap-around  $L_2$ -discrepancy. A lower bound of uniformity pattern for general asymmetric designs is provided, which can be used to serve as a benchmark for both comparing different designs and also to determine the optimal design. As a byproduct, a lower bound of wrap-around  $L_2$ -discrepancy measure for the asymmetric design is also obtained. Some illustrative examples and numerical comparisons are also provided for supporting our theoretical results.

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## 1. Introduction

Computer experiments are becoming an indispensable tool for the study of complex phenomenon or system (Santner, Williams, & Notz, 2003) and are adopted widely because of its flexibility (Wu, 2015). A lot of new challenges of modeling and designing for computer experiment arose, and space-filling property is a well known essential demand for designing computer experiment (Bates, Buck, Riccomagno, & Wynn, 1996). In case of simulation studies where a little knowledge is known about the function to be modeled, the uniform design which seeks experimental points to be uniformly scattered in the experimental domain has a wide application. Its practical success is due to its economical and flexible experimental runs to study many factors with high levels simultaneously. In this regard, mention may be made to Fang, Li, and Sudjianto (2006).

Recently, more and more attentions are paid to the space-filling property of lower dimension projection of design for computer experiment. Orthogonality or its relaxation, such as correlation, is believed to be a useful stepping stone for this purpose. He and Tang (2013) introduced strong orthogonal array which has better lower dimension uniformity than orthogonal array with same strength. Chen and Qian (2014) and Naveau, Guillou, Cooley, and Diebolt (2009) provided a class of Latin hypercube designs which can control the orthogonality or correlation for projection designs. In particular, Joseph, Gul, and Ba (2015) improved the traditional maxmin distance criterion by introducing maximum projection design and taking the projection space-filling property into account, and compared empirically it with Latin hypercube design and existing uniform design.

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For considering the projection uniformity of a design, Fang and Qin (2005) established a link among uniformity pattern, aberration and orthogonality for two-level designs. Zhang and Qin (2006) provided further results on minimum projection uniformity criterion so as to rank different two-level designs. For a high-level design, Hickernell and Liu (2002) pioneered the notion of projection discrepancy pattern. Qin, Wang, and Chatterjee (2012) suggested a uniformity pattern for  $q$ -level designs, which is generalized to asymmetric designs by Balakrishnan, Qin, and Chatterjee (2016). It should be noted that the projection uniformity measure considered in Balakrishnan et al. (2016) and Qin et al. (2012) is different from that of Fang and Qin (2005).

Computer experiment with high-level quantitative factors may be encountered often in practice. However, the uniformity pattern mentioned above is not appropriate for such a computer experiment due to the following reason. The uniformity pattern proposed in the existing works only considers whether the levels for each factor are equal or not, but the distance among them is ignored. For a function with quantitative variables, the distinct two inputs(levels) with different distances can lead to different responses, such as response surface methodology. The main objective of this paper is to study the projection uniformity pattern of designs for computer experiment with quantitative factors at general levels in the line of Fang and Qin (2005).

The paper is organized as follows. Section 2 describes in brief the notion of uniformity pattern with reference to computer experiment and the minimum projection uniformity criterion. In Section 3, we provide the derivation of a lower bound of uniformity pattern for general asymmetric case. Some illustrative examples are provided in Section 4. Finally, Section 5 gives the concluding remark.

## 2. Uniformity pattern of design for computer experiment

Uniformity is an important kind of space-filling property and is applied widely to evaluate the performance of competitive designs. Hickernell (1998) provided many modified  $L_2$ -type discrepancies based on different reproducing kernel Hilbert spaces for measuring uniformity of designs. Among these discrepancy measures, the centered  $L_2$ - and wrap-around  $L_2$ -discrepancies are found to possess good properties. But, because of the fact that the centered  $L_2$ -discrepancy is affected significantly by the center point of the experimental domain, here we consider the wrap-around  $L_2$ -discrepancy to measure the uniformity of designs.

For any positive integer  $s$ , let  $C^s = [0, 1]^s$  and  $\mathcal{H}_n = \{\mathbf{y}_1, \mathbf{y}_2, \dots, \mathbf{y}_n\}$  be a set of  $n$  points in  $C^s$ . For any  $\mathbf{u}, \mathbf{v} \in C^s$ , let  $J(\mathbf{u}, \mathbf{v}) = J_1(u_1, v_1) \times J_2(u_2, v_2) \times \dots \times J_s(u_s, v_s)$ , where the symbol “ $\times$ ” represents the Cartesian product and, for  $1 \leq i \leq s$ ,

$$J_i(u_i, v_i) = \begin{cases} [u_i, v_i], & \text{if } u_i \leq v_i; \\ [0, v_i] \cup (u_i, 1], & \text{if } u_i > v_i. \end{cases}$$

For any  $\mathbf{u}, \mathbf{v} \in C^s$ , let  $N(\mathcal{H}_n, J(\mathbf{u}, \mathbf{v}))$  and  $Vol(J(\mathbf{u}, \mathbf{v}))$  be, respectively, the cardinality of the set  $\mathcal{H}_n \cap J(\mathbf{u}, \mathbf{v})$  and the volume of the region  $J(\mathbf{u}, \mathbf{v})$ . It is to be noted that

$$Vol(J(u_i, v_i)) = \begin{cases} v_i - u_i, & \text{if } u_i \leq v_i; \\ 1 - u_i + v_i, & \text{if } u_i > v_i, \end{cases}$$

$$Vol(J(\mathbf{u}, \mathbf{v})) = \prod_{i=1}^s Vol(J(u_i, v_i)).$$

Let  $\mathcal{G}$  be the set of all possible nonempty subsets of the set  $\{1, 2, \dots, s\}$  and  $\mathcal{G}_g = \{G \in \mathcal{G} : |G| = g\}$ . For any  $G \in \mathcal{G}_g$ , let  $C^G = [0, 1]^g$  be the  $g$ -dimensional unit hypercube. According to Hickernell (1998), the wrap-around  $L_2$ -discrepancy measure for any set  $\mathcal{H}_n$  of  $n$  points from  $C^s$  is defined as

$$[WD_2(\mathcal{H}_n)]^2 = \sum_{g=1}^s \sum_{G \in \mathcal{G}_g} \iint_{C^G \times C^G} \left[ \frac{N(\mathcal{H}_n^G, J(\mathbf{u}^G, \mathbf{v}^G))}{n} - Vol(J(\mathbf{u}^G, \mathbf{v}^G)) \right]^2 d\mathbf{u}^G d\mathbf{v}^G, \tag{1}$$

where  $\mathcal{H}_n^G$  is the projection of  $\mathcal{H}_n$  on  $C^G$  and the remaining symbols are with reference to  $G$ .

It is to be noted that if one considers only uniformity of design points on  $C^G$  corresponding to some design  $G \in \mathcal{G}_g$ , then the corresponding measure of uniformity can be defined as

$$[WD_{2,G}(\mathcal{H}_n)]^2 = \iint_{C^G \times C^G} \left[ \frac{N(\mathcal{H}_n^G, J(\mathbf{u}^G, \mathbf{v}^G))}{n} - Vol(J(\mathbf{u}^G, \mathbf{v}^G)) \right]^2 d\mathbf{u}^G d\mathbf{v}^G. \tag{2}$$

Let  $U(n, q_1 q_2 \dots q_s)$  denote a class of designs with  $n$  runs and  $s$  factors each at the levels  $q_1, q_2, \dots, q_s$ , respectively. A design  $d \in U(n, q_1 q_2 \dots q_s)$  can be viewed as an  $n \times s$  array with entries of the  $j$ th,  $1 \leq j \leq s$ , column from the set  $\{0, 1, \dots, q_j - 1\}$ . Moreover, for  $1 \leq j \leq s$ , the entries of the  $j$ th column appear equally often. Here each row represents a run (point) and each column represents a factor. A typical level combination of the  $s$  factors is represented as  $\mathbf{x} = (x_1, x_2, \dots, x_s)$ , where  $x_j = 0, 1, \dots, q_j - 1$ ,  $1 \leq j \leq s$ . It is to be noted that through the mapping  $f : x_j \rightarrow y_j = (2x_j + 1)/2q_j$ , the  $n$  level combinations can be transferred into  $n$  points in  $[0, 1]^s$ .

According to [Hickernell \(1998\)](#), for any design  $d \in \mathcal{U}(n, q_1 q_2 \cdots q_s)$ , the wrap-around  $L_2$ -discrepancy measure of uniformity, as given in (1), can be expressed as

$$[WD_2(d)]^2 = -\left(\frac{4}{3}\right)^s + \frac{1}{n^2} \sum_{i_1=1}^n \sum_{i_2=1}^n \prod_{j=1}^s \left[ \frac{3}{2} - |y_{i_1 j} - y_{i_2 j}|(1 - |y_{i_1 j} - y_{i_2 j}|) \right], \tag{3}$$

where  $y_{i_1 j} = (2x_{i_1 j} + 1)/(2q_j)$ ,  $y_{i_2 j} = (2x_{i_2 j} + 1)/(2q_j)$ ,  $x_{i_1 j}$  and  $x_{i_2 j}$  are the levels of the  $j$ th factor of the design  $d$  corresponding to the  $i_1$ th and  $i_2$ th runs, respectively,  $1 \leq i_1 \neq i_2 \leq n$ ,  $1 \leq j \leq s$ .

It is noticed by many researchers that even though two designs have the same wrap-around  $L_2$ -discrepancy, they may have different uniformities when projected to a sub-design. [Fang and Qin \(2005\)](#) initiated the work on uniformity pattern to rank two-level designs. Following [Fang and Qin \(2005\)](#), we can define for any design  $d \in \mathcal{U}(n, q_1 q_2 \cdots q_s)$

$$P_g(d) = \sum_{G \in \mathcal{G}_g} [WD_{2,G}(d)]^2, \tag{4}$$

where  $[WD_{2,G}(d)]^2$  is as defined in (2). It is to be noted that  $P_g(d)$  provides a measure of overall uniformity of the design  $d$  when it is projected on to the sets belonging to  $\mathcal{G}_g$ . Moreover, it also follows from (1) and (2) that, for any design  $d \in \mathcal{U}(n, q_1 q_2 \cdots q_s)$ ,

$$[WD_2(d)]^2 = \sum_{g=1}^s P_g(d). \tag{5}$$

For any design  $d \in \mathcal{U}(n, q_1 q_2 \cdots q_s)$ , the following theorem provides an analytical expression of  $[WD_{2,G}(d)]^2$  for  $G \in \mathcal{G}_g$ , which will be the basis of rest of the paper.

**Theorem 1.** For any design  $d \in \mathcal{U}(n, q_1 q_2 \cdots q_s)$  and for any  $G \in \mathcal{G}_g$ ,  $WD_{2,G}(d)$  can be analytically expressed as

$$[WD_{2,G}(d)]^2 = -\left(\frac{1}{3}\right)^g + \frac{1}{n^2} \sum_{i_1=1}^n \sum_{i_2=1}^n \prod_{j \in G} \left[ \frac{1}{2} - |y_{i_1 j} - y_{i_2 j}|(1 - |y_{i_1 j} - y_{i_2 j}|) \right], \tag{6}$$

where the symbols are as defined in (3).

**Proof.** The proof of this theorem follows from the following two facts as well as from Eqs. (3)–(5):

(I) for any positive integer  $s$ , we can write

$$\left(\frac{4}{3}\right)^s = 1 + \sum_{i=1}^s \binom{s}{i} \left(\frac{1}{3}\right)^i;$$

(II) from the distribution law, we get

$$\begin{aligned} & \prod_{j=1}^s \left[ \frac{3}{2} - |y_{i_1 j} - y_{i_2 j}|(1 - |y_{i_1 j} - y_{i_2 j}|) \right] \\ &= \prod_{j=1}^s \left[ 1 + \left\{ \frac{1}{2} - |y_{i_1 j} - y_{i_2 j}|(1 - |y_{i_1 j} - y_{i_2 j}|) \right\} \right] \\ &= 1 + \sum_{g=1}^s \sum_{G \in \mathcal{G}_g} \prod_{j \in G} \left[ \frac{1}{2} - |y_{i_1 j} - y_{i_2 j}|(1 - |y_{i_1 j} - y_{i_2 j}|) \right]. \end{aligned}$$

For a given  $G \in \mathcal{G}_g$ , let  $d^G$  be the design whose factors are that of  $d$  just fell into  $G$ , that is,  $d^G$  is the projection design of  $d$  indexed by  $G$ . It is trivial to check

$$[WD_{2,G}(d^G)]^2 = [WD_{2,G}(d)]^2. \tag{7}$$

According to Eqs. (4) and (5), if  $g = 1$ , one can check that

$$\begin{aligned} [WD_{2,G}(d)]^2 &= [WD_{2,G}(d^G)]^2 = [WD_2(d^G)]^2 \\ &= -\frac{1}{3} + \frac{1}{n^2} \sum_{i_1=1}^n \sum_{i_2=1}^n \left[ \frac{1}{2} - |y_{i_1} - y_{i_2}|(1 - |y_{i_1} - y_{i_2}|) \right]. \end{aligned} \tag{8}$$

If  $g = 2$ , we have  $[WD_2(d^G)]^2 = \sum_{G^* \subseteq G, |G^*|=1} [WD_{2,G^*}(d^G)]^2 + [WD_{2,G}(d^G)]^2$ , where all  $[WD_{2,G^*}(d^G)]^2$ 's have an analytical formula based on Eq. (8). Thus, depending on Eqs. (4), (5) and (8),

$$\begin{aligned} [WD_{2,G}(d)]^2 &= [WD_{2,G}(d^G)]^2 = [WD_2(d^G)]^2 - \sum_{G^* \subseteq G, |G^*|=1} [WD_{2,G^*}(d^G)]^2 \\ &= -\left(\frac{1}{3}\right)^2 + \frac{1}{n^2} \sum_{i_1=1}^n \sum_{i_2=1}^n \prod_{j \in G} \left[ \frac{1}{2} - |y_{i_1j} - y_{i_2j}|(1 - |y_{i_1j} - y_{i_2j}|) \right]. \end{aligned} \tag{9}$$

Similarly, if  $g = 3$ , we have

$$[WD_2(d^G)]^2 = \sum_{G^* \subseteq G, |G^*|=1} [WD_{2,G^*}(d^G)]^2 + \sum_{G^* \subseteq G, |G^*|=2} [WD_{2,G^*}(d^G)]^2 + [WD_{2,G}(d^G)]^2,$$

where all  $[WD_{2,G^*}(d^G)]^2$ 's have an analytical formula based on Eqs. (8) and (9). Therefore,

$$\begin{aligned} [WD_{2,G}(d)]^2 &= [WD_{2,G}(d^G)]^2 = [WD_2(d^G)]^2 - \sum_{G^* \subseteq G, |G^*|=2} [WD_{2,G^*}(d^G)]^2 \\ &\quad - \sum_{G^* \subseteq G, |G^*|=1} [WD_{2,G^*}(d^G)]^2 \\ &= -\left(\frac{1}{3}\right)^2 + \frac{1}{n^2} \sum_{i_1=1}^n \sum_{i_2=1}^n \prod_{j \in G} \left[ \frac{1}{2} - |y_{i_1j} - y_{i_2j}|(1 - |y_{i_1j} - y_{i_2j}|) \right]. \end{aligned} \tag{10}$$

Apply the same method above repeatedly till  $g = s$ , the proof is completed.

**Definition 1.** For any design  $d \in \mathcal{U}(n, q_1q_2 \cdots q_s)$ , the vector  $(P_1(d), P_2(d), \dots, P_s(d))$  is called the uniformity pattern of  $d$ . Given any two designs  $d_1, d_2 \in \mathcal{U}(n, q_1q_2 \cdots q_s)$ ,  $d_1$  will be said to have less projection uniformity if, for some integer  $r$ , we have  $P_i(d_1) = P_i(d_2)$  for  $1 \leq i \leq r-1$  and  $P_r(d_1) < P_r(d_2)$ . If no design has less projection uniformity than a design  $d$ , then  $d$  will be said to have minimum projection uniformity.

Minimum projection uniformity criterion, essentially, asks for the minimization of uniformity pattern sequentially over a given class of designs. In the next section, we provide a lower bound of uniformity pattern for general asymmetrical designs.

### 3. A lower bound of uniformity pattern for asymmetrical designs

In this section we concentrate on the derivation of a lower bound of uniformity pattern for designs belonging to the class  $\mathcal{U}(n, q_1q_2 \cdots q_s)$ . In this regard the following definition will be helpful.

**Definition 2.** A multiset is a generalization of the concept of a set that, unlike a set, allows multiple instances of the multiset's elements. In other words an indexed family,  $(a_i)$ , where  $i$  is in some index-set, may define a multiset, sometimes written as  $\{a_i\}$ , in which the multiplicity of any element  $x$  is the number of indices  $i$  such that  $a_i = x$ .

For example,  $\{a, a, b\}$  and  $\{a, b\}$  are different multisets although they are the same set.

For  $1 \leq j \leq s, 1 \leq i_1 \neq i_2 \leq n$ , let us define

$$\alpha_{i_1i_2}^j = |y_{i_1j} - y_{i_2j}|(1 - |y_{i_1j} - y_{i_2j}|),$$

where  $y_{i_1j}, y_{i_2j} \in \{(2l+1)/(2q_j) : 0 \leq l \leq q_j - 1\}$ . For  $1 \leq j \leq s$ , let  $B_j$  be the set of possible values of  $\alpha_{i_1i_2}^j$  and  $A_j = \{\alpha_{i_1i_2}^j : 1 \leq i_1 \neq i_2 \leq n\}$ . Then, it is to be noted that

$$B_j = \left\{ \frac{l(q_j - l)}{q_j^2} : 0 \leq l \leq m_j \right\},$$

where  $m_j$  is some positive integer such that either  $q_j = 2m_j$  or  $q_j = 2m_j + 1$ , and that for each  $j, A_j$  is a multiset.

The following lemma will be helpful in presenting the main result of this section.

**Lemma 1.** For any design  $d \in \mathcal{U}(n, q_1q_2 \cdots q_s)$  and for  $1 \leq j \leq s$ , the multiset  $A_j$  can be expressed as

$$A_j = \bigcup_{l=0}^{m_j} A_{jl},$$

where, when  $q_j = 2m_j + 1$  with some positive integer  $m_j$ ,  $A_{jl}$  is a multiset consisting of a single element  $l(q_j - l)/q_j^2$  with multiplicity  $t_{jl}$ , here

$$t_{jl} = \begin{cases} \frac{n(n - q_j)}{q_j}, & \text{if } l = 0, \\ \frac{2n^2}{q_j}, & \text{if } l = 1, 2, \dots, m_j; \end{cases}$$

when  $q_j = 2m_j$  with some positive integer  $m_j$ ,  $A_{jl}$  is a multiset consisting of a single element  $l(q_j - l)/q_j^2$  with multiplicity  $t_{jl}$ , here

$$t_{jl} = \begin{cases} \frac{n(n - q_j)}{q_j}, & \text{if } l = 0, \\ \frac{2n^2}{q_j}, & \text{if } l = 1, 2, \dots, m_j - 1, \\ \frac{n^2}{q_j}, & \text{if } l = m_j. \end{cases}$$

**Proof.** We first consider the proof of case  $q_j = 2m_j + 1$ . For any  $j$ ,  $1 \leq j \leq s$ , here we have  $q_j = 2m_j + 1$ , where  $m_j$  is some positive integer, and for  $1 \leq i_1 \neq i_2 \leq n$ ,  $\alpha_{i_1 i_2}^j$  takes value from the set  $B_j$ . It is to be noted that the cardinality of the set  $B_j$  is  $(q_j + 1)/2$ . Moreover, the function  $\alpha_{i_1 i_2}^j$  is symmetric about the center point of the interval  $(0, 1)$ . Let  $V = \{0, 1, \dots, q_j - 1\}$ , considering a fixed number  $l$ , where  $0 \leq l \leq m_j$ , for any level  $u \in V$ , there exist an unique level  $v = u + l \pmod{q_j} \in V$  corresponding to  $u$ . Therefore, for any given element  $l (\neq 0)$  the value  $l(q_j - l)/q_j^2$  belonging to  $B_j$ , will appear  $2n^2/q_j$  times in  $A_j$ . If  $l = 0$ , the value 0 will appear  $n(n - q_j)/q_j$  times in  $A_j$ . This proves that  $A_j$  is a multiset and, for each  $l$ ,  $1 \leq l \leq m_j$ ,  $A_{jl}$  is a multiset with multiplicity  $2n^2/q_j$  and  $A_{j0}$  is a multiset with multiplicity  $n(n - q_j)/q_j$ . This completes the proof of case  $q_j = 2m_j + 1$ . The proof of case  $q_j = 2m_j$  will follow in a similar manner.

Let  $d$  be any design in  $\mathcal{U}(n, q_1 q_2 \cdots q_s)$  and  $G$  be any set belonging to  $\mathcal{G}_g$ ,  $1 \leq g \leq s$ . Let  $G_O$  and  $G_E$  be two subsets of  $G$  which are such that  $G_O \cup G_E = G$  and  $G_O \cap G_E = \emptyset$ . Moreover, suppose, for  $j \in G_O$ , there exists positive integers  $m_j$  such that  $q_j = 2m_j + 1$  and, for  $j \in G_E$ , there exists positive integers  $m_j$  such that  $q_j = 2m_j$ . Then we have the following lemma which provides a lower bound to  $[WD_{2,G}(d)]^2$ .

**Lemma 2.** If the given design  $d$  is projected on to the set  $G$ , then we have

$$[WD_{2,G}(d)]^2 \geq LB_G(d), \tag{11}$$

where

$$LB_G(d) = C_g + \frac{n-1}{n} e^S, \quad C_g = -\left(\frac{1}{3}\right)^g + \frac{1}{n} \left(\frac{1}{2}\right)^g,$$

and

$$\begin{aligned} n(n-1)S = & \sum_{j \in G} \frac{n(n - q_j)}{q_j} \cdot \ln\left(\frac{1}{2}\right) + \sum_{j \in G} \sum_{l=1}^{m_j-1} \left[ \frac{2n^2}{q_j} \cdot \ln\left(\frac{1}{2} - \frac{l(q_j - l)}{q_j^2}\right) \right] + \\ & \sum_{j \in G_O} \left[ \frac{2n^2}{q_j} \cdot \ln\left(\frac{1}{2} - \frac{m_j(m_j + 1)}{q_j^2}\right) \right] + \sum_{j \in G_E} \frac{n^2}{q_j} \cdot \ln\left(\frac{1}{4}\right). \end{aligned}$$

**Proof.** Based on Eq. (6), when the design  $d$  is projected on to  $G$ , we get

$$\begin{aligned} [WD_{2,G}(d)]^2 &= C_g + \frac{1}{n^2} \sum_{i_1=1}^n \sum_{i_2(\neq i_1)=1}^n \prod_{j \in G} \left[ \frac{1}{2} - |y_{i_1 j} - y_{i_2 j}|(1 - |y_{i_1 j} - y_{i_2 j}|) \right] \\ &= C_g + \frac{1}{n^2} \sum_{i_1=1}^n \sum_{i_2(\neq i_1)=1}^n e^{\sum_{j \in G} \ln\left(\frac{1}{2} - |y_{i_1 j} - y_{i_2 j}|(1 - |y_{i_1 j} - y_{i_2 j}|)\right)}. \end{aligned}$$

Since the function  $e^x$  is strictly convex, from Jensen's inequality, we know that

$$[WD_{2,G}(d)]^2 \geq C_g + \frac{n-1}{n} e^S,$$

where

$$S = \frac{1}{n(n-1)} \sum_{j \in G} \sum_{i_1=1}^n \sum_{i_2(\neq i_1)=1}^n \ln \left( \frac{1}{2} - |x_{ik} - y_{jk}|(1 - |x_{ik} - y_{jk}|) \right).$$

From Lemma 1, it follows that

$$\begin{aligned} n(n-1)S &= \sum_{j \in G} \frac{n(n-q_j)}{q_j} \cdot \ln \left( \frac{1}{2} \right) + \sum_{j \in G} \sum_{l=1}^{m_j-1} \left[ \frac{2n^2}{q_j} \cdot \ln \left( \frac{1}{2} - \frac{l(q_j-l)}{q_j^2} \right) \right] + \\ &\quad \sum_{j \in G_0} \left[ \frac{2n^2}{q_j} \cdot \ln \left( \frac{1}{2} - \frac{m_j(m_j+1)}{q_j^2} \right) \right] + \sum_{j \in G_E} \frac{n^2}{q_j} \cdot \ln \left( \frac{1}{4} \right). \end{aligned}$$

This completes the proof of Lemma 2.

The following theorem provides a lower to bound to the measure of uniformity pattern of any design when projected on a subdesign involving  $g$  factors.

**Theorem 2.** For a given design  $d \in \mathcal{U}(n, q_1 q_2 \cdots q_s)$  and for any  $1 \leq g \leq s$ , we have

$$P_g(d) \geq LBP_g, \tag{12}$$

where  $LBP_g = \sum_{G \in \mathcal{G}_g} LB_G(d)$  and  $LB_G(d)$  is as defined in Lemma 2.

**Proof.** Based on Eq. (4) and Lemma 2, the proof of this theorem follows immediately.

**Corollary 1.** Consider a design  $d \in \mathcal{U}(n, q^s)$  and a set  $G \in \mathcal{G}_g$ . Then we have

$$[WD_{2,G}(d)]^2 \geq LB_G(d), \tag{13}$$

where  $LB_G(d) = C_g + \left(\frac{n-1}{n}\right) e^S$ ,  $C_g = -\left(\frac{1}{3}\right)^g + \frac{1}{n} \left(\frac{1}{2}\right)^g$ , and for some positive integer  $m$ ,

$$S = \begin{cases} \frac{g(n-q)}{(n-1)q} \ln \left( \frac{1}{2} \right) + \frac{2ng}{(n-1)q} \sum_{l=1}^m \left[ \frac{1}{2} - \frac{l(q-l)}{q^2} \right], & \text{if } q = 2m + 1; f \\ \frac{g(n-q)}{(n-1)q} \ln \left( \frac{1}{2} \right) + \frac{ng}{(n-1)q} \ln \left( \frac{1}{4} \right) + \frac{2ng}{(n-1)q} \sum_{l=1}^{m-1} \left[ \frac{1}{2} - \frac{l(q-l)}{q^2} \right], & \text{if } q = 2m. \end{cases}$$

Based on Corollary 1, a theorem similar to Theorem 2 can be established related to  $P_g(d)$  values.

With reference to Eq. (5), Chatterjee, Fang, and Qin (2005) obtained the following lower bound of wrap-around  $L_2$ -discrepancy measure for a design  $d$  belonging to  $\mathcal{U}(n, 2^{m_1} 3^{m_2})$ ,

$$[WD_2(d)]^2 \geq LB_1(d), \tag{14}$$

where

$$LB_1(d) = -\left(\frac{4}{3}\right)^{m_1+m_2} + \frac{1}{n^2} \left(\frac{5}{4}\right)^{m_1} \left(\frac{23}{18}\right)^{m_2} \sum_{i=0}^{m_1} \sum_{j=0}^{m_2} \binom{m_1}{i} \binom{m_2}{j} \left(\frac{1}{5}\right)^i \left(\frac{4}{23}\right)^j \theta_{ij},$$

$\theta_{ij} = n\theta_{ij}^* + (n - 2^i 3^j \theta_{ij}^*)(\theta_{ij}^* + 1)$  and  $\theta_{ij}^*$  is the largest integer contained in  $\frac{n}{2^i 3^j}$  for  $i = 0, 1, \dots, m_1$  and  $m_2 = 0, 1, \dots, m_2$ .

Along the line of Theorem 2, we obtain a lower bound of wrap-around  $L_2$ -discrepancy, stated in Eq. (5), for general asymmetrical design. Let  $R = \{1, 2, \dots, s\}$ . Let  $R_0$  and  $R_E$  be two subsets of  $R$  which are such that  $R_0 \cup R_E = R$  and  $R_0 \cap R_E = \emptyset$ . Moreover, suppose, for  $j \in R_0$ , there exists positive integers  $m_j$  such that  $q_j = 2m_j + 1$  and, for  $j \in R_E$ , there exists positive integers  $m_j$  such that  $q_j = 2m_j$ . Then we have the following theorem which provides a lower bound to  $[WD_2(d)]^2$ .

**Theorem 3.** Given any design  $d \in \mathcal{U}(n, q_1 q_2 \cdots q_s)$ , we have

$$[WD_2(d)]^2 \geq LB_2(d), \tag{15}$$

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**Table 1**  
Two four-level designs  $d_1$  and  $d_2$  with 9 factors and 32 runs.

Runs	1	2	3	4	5	6	7	8	9	1	2	3	4	5	6	7	8	9
1	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
2	0	0	0	1	1	1	1	1	1	0	0	0	1	1	1	1	1	1
3	0	1	1	0	1	2	2	2	2	0	1	1	0	1	2	2	2	2
4	0	1	1	1	0	3	3	3	3	0	1	1	1	0	3	3	3	3
5	0	2	2	2	2	0	0	2	3	0	2	2	2	2	0	0	2	3
6	0	2	2	3	3	1	1	3	2	0	2	2	3	3	1	1	3	2
7	0	3	3	2	3	2	2	0	1	0	3	3	2	3	2	2	0	1
8	0	3	3	3	2	3	3	1	0	0	3	3	3	2	3	3	1	0
9	1	0	1	2	2	1	3	0	2	1	0	1	2	3	0	3	3	1
10	1	0	1	3	3	0	2	1	3	1	0	1	3	2	1	2	2	0
11	1	1	0	2	3	3	1	2	0	1	1	2	0	3	3	1	0	0
12	1	1	0	3	2	2	0	3	1	1	1	2	1	2	2	0	1	1
13	1	2	3	0	0	1	3	2	1	1	2	3	0	1	0	3	1	2
14	1	2	3	1	1	0	2	3	0	1	2	3	1	0	1	2	0	3
15	1	3	2	0	1	3	1	0	3	1	3	0	2	1	3	1	2	3
16	1	3	2	1	0	2	0	1	2	1	3	0	3	0	2	0	3	2
17	2	0	2	0	2	3	2	3	1	2	0	2	2	0	3	2	1	2
18	2	0	3	1	3	3	0	2	2	2	0	3	0	2	2	1	3	3
19	2	1	2	2	0	1	2	1	0	2	1	0	3	3	0	2	1	3
20	2	1	3	3	1	1	0	0	3	2	1	3	2	1	1	0	3	0
21	2	2	0	0	3	2	3	1	3	2	2	0	1	3	2	3	2	0
22	2	2	1	1	2	2	1	0	0	2	2	1	3	1	3	0	0	1
23	2	3	0	2	1	0	3	3	2	2	3	1	1	2	0	1	0	2
24	2	3	1	3	0	0	1	2	1	2	3	2	0	0	1	3	2	1
25	3	0	2	3	1	2	3	2	0	3	0	2	3	1	2	3	0	3
26	3	0	3	2	0	2	1	3	3	3	0	3	1	3	3	0	2	2
27	3	1	2	1	3	0	3	0	1	3	1	0	2	2	1	3	0	2
28	3	1	3	0	2	0	1	1	2	3	1	3	3	0	0	1	2	1
29	3	2	0	3	0	3	2	0	2	3	2	0	0	2	3	2	3	1
30	3	2	1	2	1	3	0	1	1	3	2	1	2	0	2	1	1	0
31	3	3	0	1	2	1	2	2	3	3	3	1	0	3	1	0	1	3
32	3	3	1	0	3	1	0	3	0	3	3	2	1	1	0	2	3	0

where  $LB_2(d) = C_s + \frac{n-1}{n} e^S$ ,  $C_s = -(\frac{1}{3})^s + \frac{1}{n} (\frac{1}{2})^s$  and

$$(n-1)S = \sum_{j=1}^s \frac{n-q_j}{q_j} \cdot \ln\left(\frac{3}{2}\right) + \sum_{j=1}^s \sum_{l=1}^{m_j-1} \left(\frac{2n}{q_j}\right) \ln\left[\frac{3}{2} - \frac{l(q_j-l)}{q_j^2}\right] + \sum_{j \in R_0} \left[\frac{2n}{q_j} \cdot \ln\left(\frac{3}{2} - \frac{m_j(m_j+1)}{q_j^2}\right)\right] + \sum_{j \in R_E} \frac{n}{q_j} \cdot \ln\left(\frac{5}{4}\right).$$

**4. Some illustrative examples**

In this section, some illustrative examples are provided as an application of our theoretical findings.

For designs belonging to the class  $\mathcal{U}(n, 2^{m_1} 3^{m_2})$ , we found through examples that for most of the cases the lower bound  $LB_2(= LB_2(d))$  is tighter than  $LB_1(= LB_1(d))$ . Fig. 1 displays few numerical comparisons between these two lower bounds for some values of  $m_1$  and  $m_2$ . From Fig. 1, for different values of  $m_1$ , we find that the new lower bound  $LB_2$  performs better than the existing lower bound  $LB_1$  when  $m_2$  is big and  $n$  is small, and the advantage of  $LB_1$  becomes more apparent when the ratio  $m_2/n$  goes bigger. It means our lower bound is tighter for supersaturated designs. When  $n$  is much bigger than  $m_2$ ,  $LB_1$  will be better. On this occasion, the experiment needs more runs so that the design is not economical and may not be acceptable. The similar phenomenon also occurs if we fix  $m_2$  and consider the difference  $LB_2 - LB_1$  by changing  $m_1$  and  $n$ .

**Example 1.** Consider two four-level  $U$ -type designs  $d_1$  and  $d_2$  involving 9 factors and 32 runs given in Table 1. The wrap-around  $L_2$ -discrepancy measures of  $d_1$  and  $d_2$  are 1.055390 and 1.055353, respectively. This means the design  $d_2$  is better than  $d_1$  based on wrap-around  $L_2$ -discrepancy. However, following Definition 1, it follows from Table 2 that design  $d_1$  is preferable than  $d_2$ .

**Example 2.** Two four-level designs,  $d_1 \in \mathcal{U}(8, 4^6)$  and  $d_2 \in \mathcal{U}(8, 4^{13})$  given in Table 3, are investigated to illustrate the advantage of the lower bound of uniformity pattern. The uniformity pattern and lower bounds are tabulated in Tables 4 and 5 for these two designs, respectively.

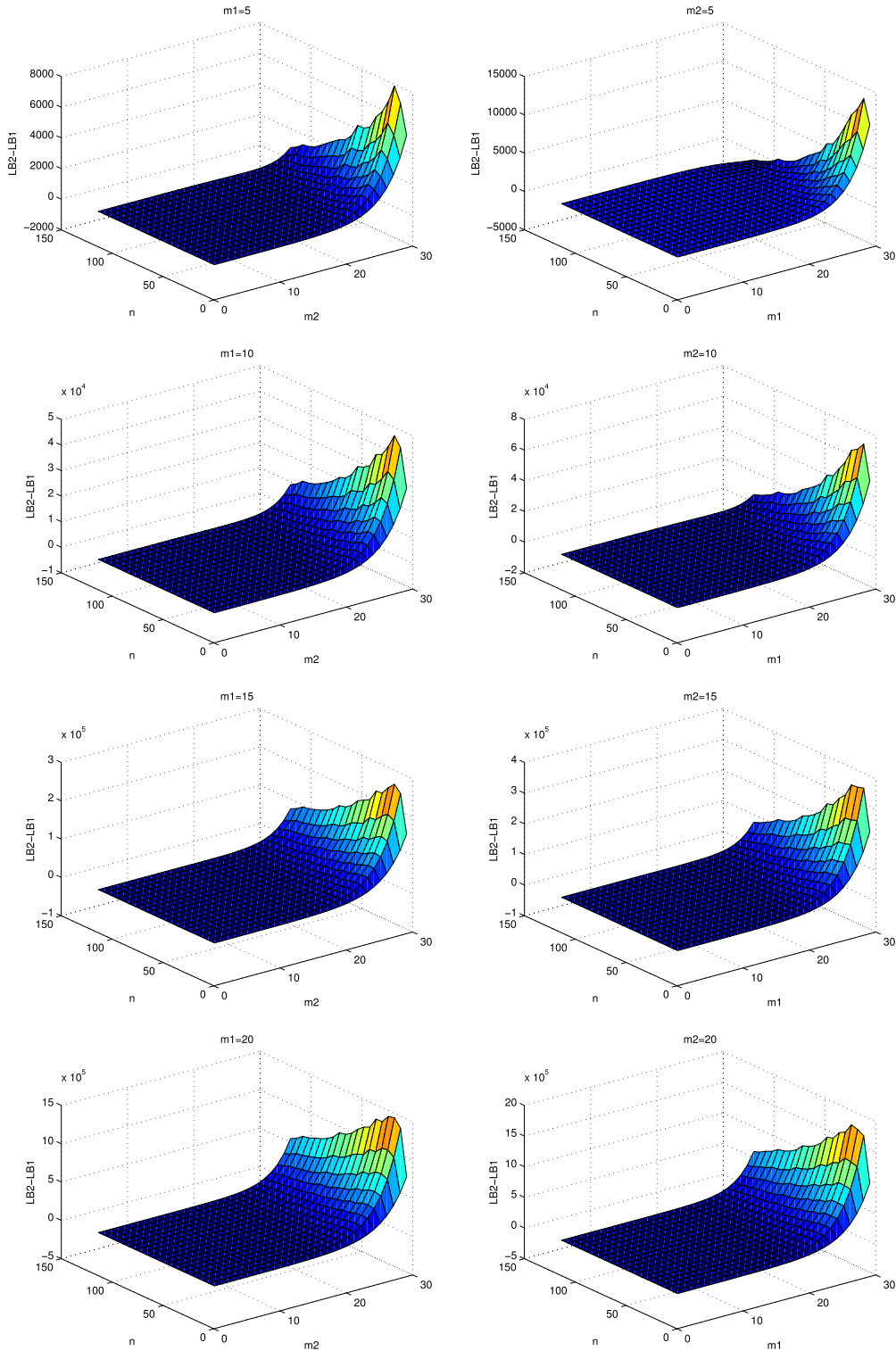


Fig. 1. Comparison between  $LB_1$  and  $LB_2$ .

**Example 3.** Consider two five-level designs with 10 runs and 9 factors, denoted by  $d_1$  and  $d_2$ . The designs  $d_1$  and  $d_2$  are provided in Table 6, and their uniformity patterns are depicted in Tables 7 and 8, respectively. Based on Tables 7 and 8, one



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**Table 2**

The uniformity pattern of design  $d_1$  and  $d_2$ .

$i$	1	2	3	4	5	6	7	8	9
$P_i(d_1)$	0.093750	0.253906	0.314006	0.233264	0.114029	0.037404	0.007975	0.001001	0.000056
$P_i(d_2)$	0.093750	0.253906	0.314010	0.233257	0.114012	0.037391	0.007970	0.001000	0.000056

**Table 3**

Two four-level designs  $d_1$  and  $d_2$  with 8 runs.

Runs	$d_1$						$d_2$												
	1	2	3	4	5	6	1	2	3	4	5	6	7	8	9	10	11	12	13
1	2	3	1	3	3	1	1	3	1	1	2	1	3	1	1	1	1	1	0
2	3	0	2	2	2	2	2	1	2	3	2	0	0	3	2	0	0	1	2
3	0	2	0	2	1	1	0	3	1	0	0	2	1	3	0	0	0	2	1
4	0	3	3	1	3	3	3	0	2	2	1	1	1	1	2	0	3	3	3
5	3	1	3	3	0	0	0	1	3	2	0	3	3	0	0	3	1	3	0
6	1	1	1	1	2	0	2	2	0	3	1	0	2	3	0	1	2	0	3
7	1	2	2	0	0	2	1	0	0	0	3	3	0	2	1	3	3	0	2
8	2	0	0	0	1	3	3	2	3	1	3	2	2	2	2	2	2	2	1

**Table 4**

The uniformity pattern and its lower bound of design  $d_1$ .

$i$	1	2	3	4	5	6
$P_i(d_1)$	0.062500	0.142904	0.134137	0.063468	0.015004	<b>0.001413</b>
$LB_i^P(d_1)$	0.021193	0.092536	0.111273	0.058880	0.014660	<b>0.001413</b>

**Table 5**

The uniformity pattern and its lower bound of design  $d_2$ .

$i$	1	2	3	4	5	6	7
$P_i(d_2)$	0.135417	0.788411	2.069700	3.253040	3.421790	2.545180	1.375960
$LB_i^P(d_2)$	0.045918	0.481189	1.591200	2.806610	3.144500	2.424750	1.338630
$i$	8	9	10	11	12	13	
$P_i(d_2)$	0.545759	0.157849	0.032501	0.004525	<b>0.000383</b>	<b>0.000015</b>	
$LB_i^P(d_2)$	0.537493	0.156567	0.032368	0.004518	<b>0.000383</b>	<b>0.000015</b>	

**Table 6**

Two five-level designs  $d_1$  and  $d_2$  with 10 runs.

Runs	$d_1$									$d_2$								
	1	2	3	4	5	6	7	8	9	1	2	3	4	5	6	7	8	9
1	0	0	0	1	2	1	2	3	4	3	2	3	4	2	2	4	4	4
2	0	1	3	2	4	3	1	1	3	4	1	2	0	4	4	2	3	1
3	1	2	3	0	1	1	3	0	2	1	3	0	1	1	2	3	4	0
4	1	3	1	1	4	0	0	2	1	0	0	1	3	3	4	1	2	
5	2	0	4	0	0	3	0	4	0	1	2	4	3	0	4	0	2	1
6	2	4	1	4	3	4	2	0	3	3	1	3	1	0	0	3	0	2
7	3	1	4	3	3	0	4	3	2	2	0	0	2	2	0	0	3	3
8	3	2	2	4	2	2	1	2	0	2	3	2	4	4	1	1	0	0
9	4	3	2	2	1	4	4	4	4	4	4	1	2	1	3	1	1	4
10	4	4	0	3	0	2	3	1	1	0	4	4	0	3	1	2	2	3

**Table 7**

The uniformity pattern and its lower bound of design  $d_1$ .

$i$	1	2	3	4	5	6	7	8	9
$P_i(d_1)$	0.060000	0.242240	0.420441	0.409698	0.245879	0.0933912	0.0219954	0.00294492	<b>0.000171988</b>
$LB_i^P(d_1)$	0.001490	0.114864	0.301804	0.348417	0.226919	0.0898777	0.0216344	0.00292904	<b>0.000171988</b>

can find that  $d_1$  possesses less  $P_i(\cdot)$  than  $d_2$  for  $i = 1, 2, \dots, 9$ . Depending on Definition 1,  $d_1$  is better than  $d_2$ . Moreover, one component of uniformity pattern of  $d_1$  achieves the lower bound provided in Theorem 2.

**5. Conclusion and further discussion**

Computer experiment (or simulation) becomes more and more important to investigate complicated systems, and space-filling property of design is an essential demand for such experiment. However, different designs with same space-filling

**Table 8**The uniformity pattern and its lower bound of design  $d_2$ .

$i$	1	2	3	4	5	6	7	8	9
$P_i(d_2)$	0.060000	0.251584	0.440560	0.428227	0.255339	0.0962831	0.0225247	0.00299860	0.000174316
$LB_i^p(d_2)$	0.001490	0.114864	0.301804	0.348417	0.226919	0.0898777	0.0216344	0.00292904	0.000171988

property may possess different space-filling properties for lower dimension projection designs. Recently, more and more researchers are focusing on exploring designs which have good space-filling property of lower dimension projection, such as He and Tang (2013). And some new criteria are introduced to provide more abundant information of designs when lower dimension projection is considered, such as the work of Joseph et al. (2015) and Zhang and Qin (2006). In this paper, minimum projection uniformity designs are introduced and studied for computer experiment with quantitative factors. And we provide a lower bound of uniformity pattern for general symmetrical and asymmetrical designs, which is more appropriate to systematically rank and select good design for computer experiment.

We also note that a challenge problem, which refers to construct minimum projection design, is worthy studying further. The major challenge in constructing minimum projection designs lies in that this criterion asks for sequentially minimizing the uniformity pattern. Thus, it is a constraint sequential optimization problem. We should first optimize the first component of uniformity pattern, and keep its optimal value as a constraint, then we can optimize the second component, and so on. The difficulty is that when we are optimizing some component, the optimal components before it may become bad. One feasible way is to optimize the weighted sum of uniformity pattern with an appropriate weight, as shown in Xu (2005), and apply heuristic threshold accepting algorithm to avoid local minimizer. Moreover, a sharp lower bound can also be used to enhance the effect for searching minimum projection uniformity designs.

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### Appendix A. Supplementary data

Supplementary material related to this article can be found online at <https://doi.org/10.1016/j.jkss.2018.01.002>.

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